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LAST NAME:

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Solution

**Problem 1** Let  $L_1$  be the language accepted by the finite automaton given on Figure 1.

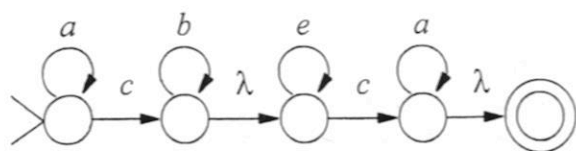


Figure 1: Language  $L_1$

Let  $L_2$  be the language generated by the context-free grammar  $G_2 = (V, \Sigma, P, S_2)$ , where  $\Sigma = \{a, b, c, e\}$ ,  $V = \{S_2\}$ , and the production set  $P$  is:

$$S_2 \rightarrow aS_2a \mid bS_2b \mid cS_2c \mid \lambda$$

(a) Write 5 distinct strings that belong to  $L_1$  but do not belong to  $L_2$  (belong to  $L_1 \setminus L_2$ ). If such strings do not exist, state it and explain why.

Answer:

acc, acbc, cca,  
decaa

(b) Write 5 distinct strings that belong to  $L_2$  but do not belong to  $L_1$  (belong to  $L_2 \setminus L_1$ ). If such strings do not exist, state it and explain why.

Answer:

$\lambda$ , aa, bb, abba, baab

(c) Write 5 distinct strings that belong to  $L_1$  and  $L_2$  (belong to  $L_1 \cap L_2$ ). If such strings do not exist, state it and explain why.

Answer:

cc, acca, aaccaa,  
cbbe, acbbca

(d) Write 5 distinct strings over alphabet  $\{a, b, c, d\}$  that do not belong to  $L_1$  and do not belong to  $L_2$  (belong to  $\overline{L_1 \cap L_2}$ ). If such strings do not exist, state it and explain why.

Answer:

ccc, c, bcb, ba, cb

(e) Write 5 distinct strings that belong to  $a^*c^*b^*e^*c^*a^*$  but do not belong to  $L_1$  (belong to  $a^*c^*b^*e^*c^*a^* \setminus L_1$ ). If such strings do not exist, state it and explain why.

Answer:

$\lambda$ , c, ccc, ab, a

(f) Write 5 distinct strings that belong to  $L_1$  but have a length equal to 3. If such strings do not exist, state it and explain why.

Answer:

only 4 exist:  
acc, cbc, ccc, cca

**Problem 3** (a) Calculate the image of the sequence  $\langle 3, 0, 1, 1 \rangle$  under Gödel numbering and show your work. If this image does not exist, state it and explain why.

Answer:

$$\begin{aligned} g(\langle 3, 0, 1, 1 \rangle) &= \\ 2^{3+1} \cdot 3^{0+1} \cdot 5^{1+1} \cdot 7^{1+1} &= \\ 16 \cdot 3 \cdot 25 \cdot 49 &= \\ 4 \cdot 100 \cdot 3 \cdot 49 &= \\ 400 \cdot 147 &= \\ \boxed{158800} \end{aligned}$$

(b) Calculate the pre-image (original) of the number 5880 under Gödel numbering and show your work. If this pre-image does not exist, state it and explain why.

Answer:

$$\begin{aligned} 5880 &= 2 \cdot 2940 \\ &= 2 \cdot 2 \cdot 1470 \\ &= 2 \cdot 2 \cdot 2 \cdot 735 \\ &= 2 \cdot 2 \cdot 2 \cdot 3 \cdot 245 \\ &= 2 \cdot 2 \cdot 2 \cdot 3 \cdot 5 \cdot 49 \\ &= 2^3 \cdot 3 \cdot 5 \cdot 7^2 \end{aligned}$$

$$\begin{aligned} g^{-1}(5880) &= \\ \boxed{\langle 2, 0, 0, 1 \rangle} \end{aligned}$$

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Soludion

In each of the cases below, state the cardinality of the given set. If this cardinality is finite, state the exact number; if it is infinite, specify whether it is countable or uncountable.

(c) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset \cup a$$

Answer:

1

(d) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset^* \cup a^*$$

Answer:

infinite and countable

(e) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset^* \cup a$$

Answer:

2

(f) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset^* a$$

Answer:

1

(g) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset a$$

Answer:

0

(h) set whose regular expression over  $\Sigma = \{a, b\}$  is:

$$\emptyset^* \cup \lambda$$

Answer:

1

(i) class of languages over  $\Sigma = \{a, b\}$  that are regular;

Answer:

infinite and countable

**Problem 5** Let  $L$  be the set of all strings over the alphabet  $\{a, b, c\}$  which satisfy all of the following conditions:

1. does not begin with  $c$ ;
2. does not end with  $a$ ;
3. has an odd length.

(a) Write 5 distinct strings that belong to  $L$ . If such strings do not exist, state it and explain why.

$b, bbb, abc, bcb,$   
 $aab$

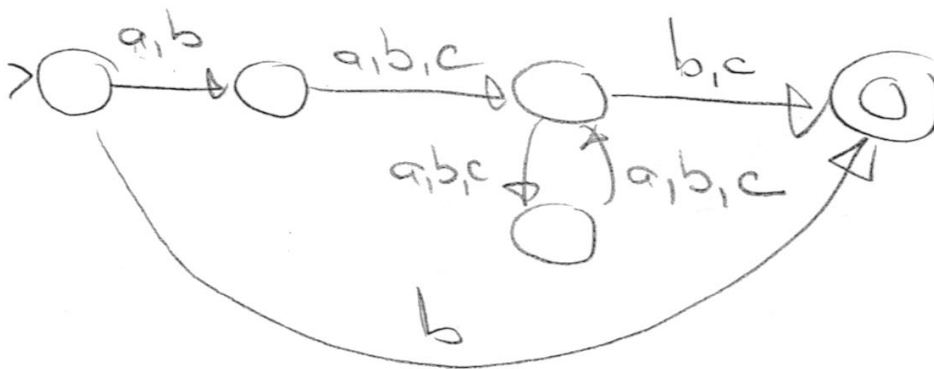
(b) Write a regular expression that represents the language  $L$ . If such a regular expression does not exist, state it and explain why.

Answer:

$b \cup (aub)(aubuc)(aubuc)(aubuc)^*(buc)$

(c) Draw a state-transition graph of a finite automaton that accepts the language  $L$ . If such an automaton does not exist, state it and explain why.

Answer:



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(d) Write a complete formal definition of a context-free grammar that generates the language  $L$ . If such a grammar does not exist, state it and explain why.

Answer:

$G = (V, \Sigma, P, S), \Sigma = \{a, b, c\}$   
 $V = \{S, A, B, Z, E\}$

$P: S \rightarrow b \mid AZEB$

$A \rightarrow a \mid b$

$B \rightarrow b \mid c$

$Z \rightarrow a \mid b \mid c$

$E \rightarrow \Lambda \mid EE \mid ZZ$

**Problem 7** Let  $L$  be the set of all strings over the alphabet  $\{a, b, c\}$  which satisfy all of the following properties.

1. if the string does not contain any  $b$ 's then its length is even and it contains exactly one  $a$ ;
2. if the string contains a positive even number of  $b$ 's, then its length is odd and it contains exactly one  $c$ ;
3. if the string contains an odd number of  $b$ 's, then it is a palindrome.

Write a complete formal definition of a context-free grammar that generates the language  $L$ . If such a grammar does not exist, state it and explain why.

**Answer:**

$$G = (V, \Sigma, P, S), \quad \Sigma = \{a, b, c\}$$

$$V = \{S, S_1, S_2, S_3, E, N_{00}, N_{01}, N_{02}, Y_{00}, Y_{01}, Y_{02}, N_{10}, N_{11}, N_{12}, Y_{10}, Y_{11}, Y_{12}\}$$

$P:$

$$S \rightarrow S_1 S_2 S_3$$

$$S_1 \rightarrow E a c E \mid E c a E$$

$$E \rightarrow c c E \mid \lambda$$

$$S_3 \rightarrow a S_3 a \mid b S_3 b \mid c S_3 c \mid \lambda$$

$$S_2 \rightarrow N_{00} \mid Y_{02} \rightarrow \lambda$$

$$N_{00} \rightarrow a N_{10} \mid b N_{01} \mid c Y_{00}$$

$$N_{01} \rightarrow a N_{11} \mid b N_{02} \mid c Y_{01}$$

$$N_{02} \rightarrow a N_{12} \mid b N_{01} \mid c Y_{02}$$

$$Y_{00} \rightarrow a Y_{10} \mid b Y_{01}$$

$$Y_{01} \rightarrow a Y_{11} \mid b Y_{02}$$

$$Y_{02} \rightarrow a Y_{12} \mid b Y_{01}$$

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$N$  : no  $c$   
 $Y$  : seen  $c$

first index : count  $a$   
 second index : count  $b$

$$N_{10} \rightarrow a N_{00} \mid b N_{11} \mid c Y_{10}$$

$$N_{11} \rightarrow a N_{01} \mid b N_{12} \mid c Y_{11}$$

$$N_{12} \rightarrow a N_{02} \mid b N_{11} \mid c Y_{12}$$

$$Y_{10} \rightarrow a Y_{00} \mid b Y_{11}$$

$$Y_{11} \rightarrow a Y_{01} \mid b Y_{12}$$

$$Y_{12} \rightarrow a Y_{02} \mid b Y_{11}$$